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RFC 9933

Carrying SR-Algorithm in Path Computation Element Communication Protocol (PCEP)

Abstract

This document specifies extensions to the Path Computation Element Communication Protocol (PCEP) to enhance support for Segment Routing (SR) with a focus on the use of Segment Identifiers (SIDs) and SR-Algorithms in Traffic Engineering (TE). The SR-Algorithm associated with a SID defines the path computation algorithm used by Interior Gateway Protocols (IGPs). It introduces mechanisms for PCEP peers to signal the SR-Algorithm associated with SIDs by encoding this information in Explicit Route Object (ERO) and Record Route Object (RRO) subobjects, enables SR-Algorithm constraints for path computation, and defines new metric types for the METRIC object. This document updates RFC 8664 and RFC 9603 to allow such extensions.

Status of This Memo

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1. Introduction

[RFC5440] describes the Path Computation Element Communication Protocol (PCEP) for communication between a Path Computation Client (PCC) and a Path Computation Element (PCE) or between a pair of PCEs. [RFC8664] and [RFC9603] specify PCEP extensions to support Segment Routing (SR) over MPLS and IPv6 data planes, respectively.

This document specifies extensions to PCEP to enhance support for SR Traffic Engineering (TE). Specifically, it focuses on the use of Segment Identifiers (SIDs) and SR-Algorithms. An SR-Algorithm associated with a SID defines the path computation algorithm used by IGP.

The PCEP extensions specified in this document are as follows:

Signaling SR-Algorithm in ERO and RRO: Mechanisms are introduced for PCEP peers to exchange information about the SR-Algorithm associated with each SID. This includes extending SR-ERO, SR-RRO, SRv6-ERO, and SRv6-RRO subobjects to carry an Algorithm field. This document updates [\[RFC8664\]](#) and [\[RFC9603\]](#) to enable such encoding.

SR-Algorithm Constraint for Path Computation: Mechanisms are defined for signaling a specific SR-Algorithm as a constraint to the PCE for path computation. This includes a new SR-Algorithm TLV carried in the Label Switched Path Attributes (LSPA) object.

Extensions to METRIC object: Several new metric types are introduced for the METRIC object to support optimization metrics derived from Flexible Algorithm Definitions (FADs) during Flexible Algorithm path computation; their application is not restricted to Flexible Algorithms, and they may be used with Label Switched Paths (LSPs) set up using different Path Setup Types (PSTs).

1.1. Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "NOT RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14 [\[RFC2119\]](#) [\[RFC8174\]](#) when, and only when, they appear in all capitals, as shown here.

2. Terminology

This document uses the following terms defined in [\[RFC5440\]](#): Explicit Route Object (ERO), Label Switched Path Attributes (LSPA), Path Computation Client (PCC), Path Computation Element (PCE), Path Computation Element Communication Protocol (PCEP), PCEP peer, PCEP speaker, Record Route Object (RRO), and Traffic Engineering Database (TED).

This document uses the following term defined in [\[RFC3031\]](#): Label Switched Path (LSP).

This document uses the following term defined in [\[RFC9479\]](#) and [\[RFC9492\]](#): Application-Specific Link Attributes (ASLA).

This document uses the following terms defined in [\[RFC8664\]](#): Node or Adjacency Identifier (NAI) and Segment Routing Database (SR-DB).

This document uses the following terms defined in [\[RFC9350\]](#): Flexible Algorithm Definition (FAD) and winning FAD.

Note that the base PCEP specification [RFC4655] originally defined the use of the PCE architecture for MPLS and GMPLS networks with LSPs instantiated using the RSVP-TE signaling protocol. Over time, support for additional PSTs, such as SRv6, has been introduced [RFC9603]. The term "LSP" is used extensively in PCEP specifications and, in the context of this document, refers to a Candidate Path within an SR Policy, which may be an SRv6 path (still represented using the LSP object as specified in [RFC8231]).

The term "extension block" is used in this document to identify the additional bytes appended to a PCEP object, which may exist depending on the inclusion of a flag in that object

The following terminologies are used in this document:

P2MP: Point-to-Multipoint

Subobject Extension Block: Optional, variable-length extension block for SR-ERO and SR-RRO subobjects defined in Section 4.2.1 of this document.

Subobject Extension Block Flag (SEBF): Any flag in the Flags field of SR-ERO or SR-RRO subobjects that is used to signal that the corresponding field is encoded in the Subobject Extension Block.

3. Motivation

Existing PCEP specifications lack mechanisms to explicitly signal and negotiate SR-Algorithm capabilities and constraints. This limits the ability of PCEs to make informed path computation decisions based on the specific SR-Algorithms supported and desired within the network. The absence of an explicit SR-Algorithm specification in PCEP messages implied no specific constraint on the SR-Algorithm to be used for path computation, effectively allowing the use of SIDs with any SR-Algorithm.

A primary motivation for these extensions is to enable the PCE to leverage the path computation logic and topological information derived from IGP, including Flexible Algorithms. Aligning PCE path computation with these IGP algorithms enables network operators to obtain paths that are congruent with the underlying routing behavior, which can result in segment lists with a reduced number of SIDs. The support for SR-Algorithm constraints in PCE path computation simplifies the deployment and management of Flexible Algorithm paths in multi-domain network scenarios.

The PCE and the PCC may independently compute SR-TE paths with different SR-Algorithms. This information needs to be exchanged between PCEP peers for purposes such as network monitoring and troubleshooting. In scenarios involving multiple PCEs, when a PCC receives a path from the primary PCE, it needs to be able to report the complete path information, including the SR-Algorithm, to a backup PCE. This is essential for high availability (HA) scenarios, ensuring that the backup PCE can correctly verify Prefix SIDs.

The introduction of an SR-Algorithm TLV within the LSPA object allows operators to specify SR-Algorithm constraints directly, thereby refining path computations to meet specific needs, such as low-latency paths.

The ability to specify an SR-Algorithm per SID in ERO and RRO is crucial for multiple reasons, for example:

- SID types without algorithm specified - Certain SID types, such as Binding SIDs (BSIDs) [RFC8402], may not have an SR-Algorithm specified. It may be inaccurate to state that an entire end-to-end path adheres to a specific algorithm if it includes a BSID from another policy.

Note: In SRv6, the BSID can be allocated from an algorithm-specific SRv6 Locator, which will result in the path to that BSID PCC node following that algorithm-specific path. However, the implicit algorithm of BSID is independent of the SR-Algorithm used for the SR Policy associated with that BSID.

- Topologies with two IGP domains, each using the same FAD but with differing algorithm numbers.

4. Object Formats

4.1. OPEN Object

4.1.1. SR PCE Capability Sub-TLV

The SR-PCE-CAPABILITY sub-TLV is defined in [Section 4.1.2](#) of [RFC8664] to be included in the PATH-SETUP-TYPE-CAPABILITY TLV.

This document defines the following flag in the SR-PCE-CAPABILITY Sub-TLV Flags field:

S (SR-Algorithm Capability) - bit 5: If the S flag is set to 1, a PCEP speaker indicates support for the Algorithm field and the Subobject Extension Block in the SR-ERO subobject described in [Section 4.2](#) and the SR-Algorithm TLV described in [Section 4.4](#) for LSPs set up using PST 1 (Segment Routing) [RFC8664]. It does not indicate support for these extensions for other PSTs. If the S flag is set to 0, behavior reverts to the procedures defined in existing specifications prior to the introduction of this extension.

4.1.2. SRv6 PCE Capability Sub-TLV

The SRv6-PCE-CAPABILITY sub-TLV is defined in [Section 4.1.1](#) of [RFC9603] to be included in the PATH-SETUP-TYPE-CAPABILITY TLV.

This document defines the following flag in the SRv6-PCE-CAPABILITY Sub-TLV Flags field:

SR-Algorithm Capability (S) - bit 13:

If the S flag is set to 1, a PCEP speaker indicates support for the Algorithm field in the SRv6-ERO subobject described in [Section 4.3](#) and the SR-Algorithm TLV described in [Section 4.4](#) for LSPs set up using PST 3 (SRv6) [[RFC9603](#)]. It does not indicate support for these extensions for other PSTs. If the S flag is set to 0, behavior reverts to the procedures defined in existing specifications prior to the introduction of this extension.

4.2. SR-ERO Subobject

This document updates the SR-ERO subobject format defined in [Section 4.3.1](#) of [[RFC8664](#)] with a new optional, variable-length Subobject Extension Block field. The block is used to convey additional information, such as the Algorithm field, and is designed to allow future extensibility. Further, a new A flag in the Flags field is introduced as shown in [Figure 1](#).

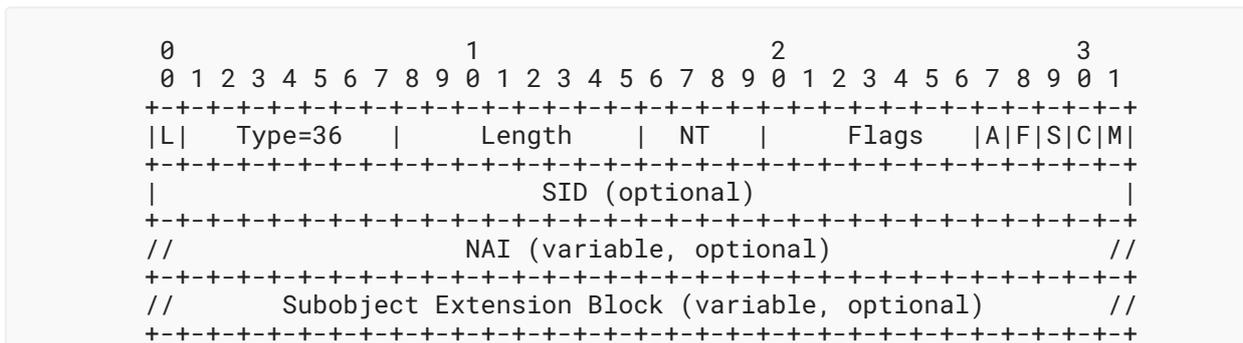


Figure 1: SR-ERO Subobject Format

A new flag in the Flags field:

A (SR-Algorithm Flag): If set by a PCEP speaker, the Subobject Extension Block **MUST** be included in the SR-ERO subobject, as shown in [Figure 1](#), along with the specified algorithm. The length of this block is variable and determined by subtracting the size of the fixed fields and any optional SID or NAI fields from the total subobject Length. The length of the Subobject Extension Block **MUST** always be a multiple of 4 bytes. If this flag is set to 0, then either:

- the Subobject Extension Block is not included and processing described in [Section 5.2.1](#) of [[RFC8664](#)] applies or
- the Subobject Extension Block is included (due to an SEBF in a future specifications) and the Algorithm field **MUST** be ignored.

This document updates the SR-ERO subobject validation defined in [Section 5.2.1](#) of [[RFC8664](#)] by extending existing validation to include the Subobject Extension Block and the A flag, as follows.

On receiving an SR-ERO subobject, a PCC **MUST** validate that the Length field, S bit, F bit, A bit, NT field, and any present SEBFs are consistent, as follows:

- If the Subobject Extension Block is included (i.e., if any SEBF, such as A or a future flag, is set to 1), the length of the subobject **MUST** include the size of the entire Subobject Extension Block as determined by the set of SEBFs.
 - The minimum size of the Subobject Extension Block is 4 bytes when only a single SEBF (such as A) is set and may be longer (in multiples of 4 bytes) if additional SEBFs are set and require more space.
 - The total subobject Length is the sum of the sizes of the fixed and optional fields (SID, NAI, etc.) and the total size of the Subobject Extension Block required by the set of SEBFs.
 - The exact calculation of Length for each NT, S, F, and set of SEBFs is as follows:
 - If NT=0, the F bit **MUST** be 1, the S bit **MUST** be 0, and the Length **MUST** be 8 + the size of the Subobject Extension Block.
 - If NT=1, the F bit **MUST** be 0.
 - If the S bit is 1, the Length **MUST** be 8 + the size of the Subobject Extension Block.
 - If the S bit is 0, the Length **MUST** be 12 + the size of the Subobject Extension Block.
 - If NT=2, the F bit **MUST** be 0.
 - If the S bit is 1, the Length **MUST** be 20 + the size of the Subobject Extension Block.
 - If the S bit is 0, the Length **MUST** be 24 + the size of the Subobject Extension Block.
 - If NT=3, the F bit **MUST** be 0.
 - If the S bit is 1, the Length **MUST** be 12 + the size of the Subobject Extension Block.
 - If the S bit is 0, the Length **MUST** be 16 + the size of the Subobject Extension Block.
 - If NT=4, the F bit **MUST** be 0.
 - If the S bit is 1, the Length **MUST** be 36 + the size of the Subobject Extension Block.
 - If the S bit is 0, the Length **MUST** be 40 + the size of the Subobject Extension Block.
 - If NT=5, the F bit **MUST** be 0.
 - If the S bit is 1, the Length **MUST** be 20 + the size of the Subobject Extension Block.
 - If the S bit is 0, the Length **MUST** be 24 + the size of the Subobject Extension Block.
 - If NT=6, the F bit **MUST** be 0.
 - If the S bit is 1, the Length **MUST** be 44 + the size of the Subobject Extension Block.
 - If the S bit is 0, the Length **MUST** be 48 + the size of the Subobject Extension Block.
- If no SEBF (including the A flag defined in this document) is set, the Length value **MUST** follow the requirements defined in [Section 5.2.1](#) of [RFC8664].

4.2.1. Subobject Extension Block

The Subobject Extension Block is an optional, extensible field in the SR-ERO subobject. Its presence is indicated by the setting of any SEBF in the subobject's Flags field (e.g., the A flag defined in this document or flags defined by future specifications).

Block length and presence:

- If the A flag is set, and no other SEBF is set, the block length **MUST** be 4.
- The block length is at least 4 bytes when present.
- The block length **MUST** always be a multiple of 4 bytes.
- The block **MUST** be included if any SEBF is set in the Flags field.
- Future extensions may define additional SEBFs and corresponding fields, allowing the block to be increased in size beyond the initial 4 bytes as needed.

The first 4 bytes of the Subobject Extension Block are described in [Figure 2](#).



Figure 2: Subobject Extension Block Format

Unassigned (24 bits):

This field is reserved for future use and **MUST** be set to zero when sending and ignored when receiving.

Algorithm (8 bits):

The SR-Algorithm value from the "IGP Algorithm Types" registry of the "Interior Gateway Protocol (IGP) Parameters" registry group (see [IANA-ALGORITHM-TYPES](#)).

Future extensions **SHOULD** first use the Unassigned portion of the initial 4 bytes to carry new information. If additional space is needed, the Subobject Extension Block may be extended in 4-byte increments. Each such extension must be indicated by a dedicated SEBF in the Flags field (similar to the A flag) and must be accompanied by capability signaling in an appropriate capability sub-TLV. The specific sub-TLV to be used is not restricted by this specification and may include, for example, the SR-PCE-CAPABILITY sub-TLV, the SRv6-PCE-CAPABILITY sub-TLV, or other capability TLVs, depending on the context of the extension. Interoperability procedures and the precise signaling mechanisms for each new SEBF and its associated capability will be defined by future specifications or procedures describing those extensions.

When receiving a Subobject Extension Block longer than 4 bytes, receivers that do not recognize or have not negotiated support for additional flags **MUST** ignore the unknown additional bytes beyond those defined in this document.

4.2.2. Guidance for Future Extensions

Future enhancements extending the Subobject Extension Block must:

- Define a new SEBF in the Flags field to indicate the presence of a new extension and specify the corresponding capability signaling for that extension.
- Specify which parts of the reserved/extension block are used and how the block length is calculated when their extension is present.
- Ensure the reserved bits in the initial 4 bytes are used when possible and the block is extended only when additional space is necessary.
- Have future extensions define additional SEBFs and corresponding fields, allowing the block to be increased in size beyond the initial 4 bytes as needed.

Example: Future extension introducing a Z flag and a new Z field (8 bits):

- If the A flag and/or the Z flag are set, the Subobject Extension Block is included. The Z field may use 8 bits of the reserved portion. A field is only considered valid if its corresponding flag is set. For example, if the Z flag is set to 1 but the A flag is set to 0, the Z field is valid but the Algorithm field is ignored.
- If space beyond the initial 4 bytes is needed, the extension document specifies the new block layout and total length. To simplify parsing, if a flag for such an extension is set, the full extended block is encoded, including the initial 4 bytes, even if the A flag is set to 0.

4.3. SRv6-ERO Subobject

This document updates the SRv6-ERO subobject format defined in Section 4.3.1 of [RFC9603] with the Algorithm field carved out of the Reserved field. Further, a new A flag is defined in the existing Flags field as shown in Figure 3.

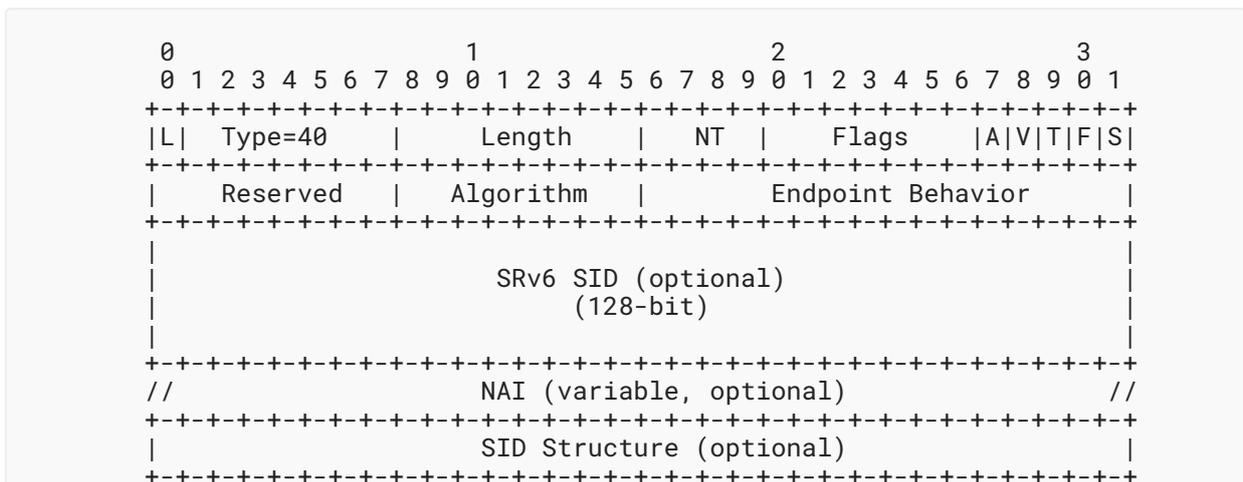


Figure 3: SRv6-ERO Subobject Format

Flags field:

A (SR-Algorithm Flag): If set by a PCEP speaker, the Algorithm field is included in the SRv6-ERO subobject as specified in [Figure 3](#). If this flag is set to 0, then the Algorithm field is absent and processing described in [Section 5.2.1](#) of [\[RFC9603\]](#) applies.

Reserved (8 bits):

It **MUST** be set to 0 while sending and ignored on receipt.

Algorithm (8 bits):

The SR-Algorithm value from the "IGP Algorithm Types" registry of the "Interior Gateway Protocol (IGP) Parameters" registry group.

Note: The Subobject Extension Block is applicable to the SRv6-ERO subobject but is not required by this specific specification as existing reserved space is used. When additional space is needed in the SRv6-ERO subobject, the future extensions **SHOULD** specify the usage of the Subobject Extension Block for the SRv6-ERO subobject.

4.4. SR-Algorithm TLV

A new TLV for the LSPA object is introduced to carry the SR-Algorithm constraint ([Section 5.2](#)). This TLV **MUST** only be used when PST = 1 or 3 for SR-MPLS and SRv6, respectively. Only the first instance of this TLV **MUST** be processed; subsequent instances **MUST** be ignored.

The format of the SR-Algorithm TLV is as follows:

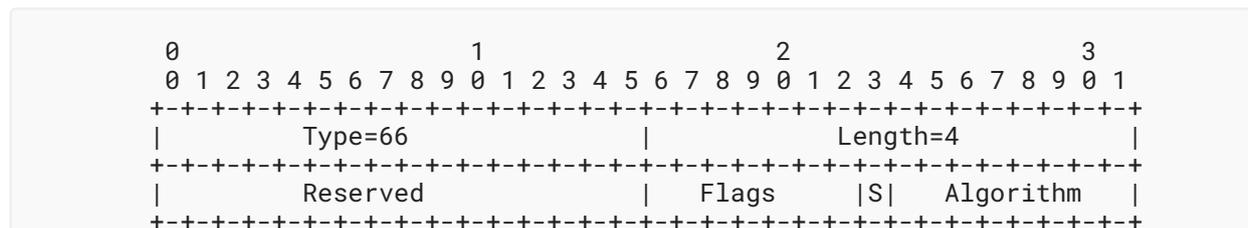


Figure 4: SR-Algorithm TLV Format

Type (16 bits): 66

Length (16 bits): 4

The 32-bit value is formatted as follows.

Reserved (16 bits): **MUST** be set to 0 by the sender and **MUST** be ignored by the receiver.

Flags (8 bits): This document defines the following flag. The other flags **MUST** be set to 0 by the sender and **MUST** be ignored by the receiver.

S (Strict):

If set, the path computation at the PCE **MUST** fail if the specified SR-Algorithm constraint cannot be satisfied. If the S (Strict) bit is unset and the PCE is unable to compute a path that satisfies the specified SR-Algorithm constraint, the PCE **MUST** attempt to compute a path as if no SR-Algorithm constraint had been requested. This means the PCE may use any available SR-Algorithm for the computation, consistent with the default behavior in the absence of SR-Algorithm constraint.

Algorithm (8 bits): The SR-Algorithm to be used during path computation (see [Section 5.2](#)).

4.5. Extensions to METRIC Object

The METRIC object is defined in [Section 7.8](#) of [\[RFC5440\]](#). This document specifies additional types for the METRIC object to enable the encoding of optimization metric types derived from the FAD during Flexible Algorithm path computation (see [Section 5.2.2](#)). While these new metric types are defined to support this specific use case, their use is not restricted to Flexible Algorithm path computation or to any specific PST.

- T=22: Path Min Delay metric ([Section 4.5.1.1](#))
- T=23: P2MP Path Min Delay metric ([Section 4.5.1.2](#))
- T=24: Path Bandwidth metric ([Section 4.5.2.1](#))
- T=25: P2MP Path Bandwidth metric ([Section 4.5.2.2](#))
- T=128-255: User-defined metric ([Section 4.5.3](#))

The following terminology is used and expanded along the way.

- A network comprises a set of N links $\{L_i, (i=1\dots N)\}$.
- A path P of a point-to-point (P2P) LSP is a list of K links $\{L_{pi}, (i=1\dots K)\}$.
- A P2MP tree T comprises a set of M destinations $\{Dest_j, (j=1\dots M)\}$.

4.5.1. Path Min Delay Metric

[\[RFC7471\]](#) and [\[RFC8570\]](#) define the "Min/Max Unidirectional Link Delay" sub-TLV to advertise the link minimum and maximum delay in microseconds in a 24-bit field.

[\[RFC5440\]](#) defines the METRIC object with a 32-bit metric value encoded in IEEE floating point format (see [\[IEEE.754.2019\]](#)).

The encoding for the Path Min Delay metric value is quantified in units of microseconds and encoded in IEEE floating point format.

For use in the PCEP METRIC object, the 24-bit unsigned integer delay value is converted to a 32-bit IEEE floating point value. This conversion follows the procedure specified in [\[IEEE.754.2019\]](#).

4.5.1.1. P2P Path Min Delay Metric

The minimum Link Delay metric is defined in [\[RFC7471\]](#) and [\[RFC8570\]](#) as "Min Unidirectional Link Delay". The Path Min Link Delay metric represents the measured minimum link delay value over a configurable interval.

The Path Min Delay metric type of the METRIC object in PCEP represents the sum of the Min Link Delay metric of all links along a P2P path.

- A Min Link Delay metric of link L is denoted by $D(L)$.
- A Path Min Delay metric for the P2P path $P = \text{Sum} \{D(L_{pi}), (i=1\dots K)\}$.

4.5.1.2. P2MP Path Min Delay Metric

The P2MP Path Min Delay metric type of the METRIC object in PCEP encodes the Path Min Delay metric for the destination that observes the worst (i.e., highest value) delay metric among all destinations of the P2MP tree.

- The P2P Path Min Delay metric of the path to destination Dest_j is denoted by $\text{PMDM}(\text{Dest}_j)$.
- The P2MP Path Min Delay metric for the P2MP tree $T = \text{Maximum}\{\text{PMDM}(\text{Dest}_j), (j=1\dots M)\}$.

4.5.2. Path Bandwidth Metric

Section 4 of [RFC9843] defines a new metric type, "Bandwidth metric", which may be advertised in their link metric advertisements.

When performing Flexible Algorithm path computation as described in Section 5.2.2, procedures described in Sections 4.1 and 5 from [RFC9843] **MUST** be followed with automatic metric calculation.

For path computations in contexts other than Flexible Algorithm (including PSTs other than 1 or 3 for SR-MPLS and SRv6, respectively), if the Generic Metric sub-TLV with the Bandwidth metric type is not advertised for a link, the PCE implementation **MAY** apply a local policy to derive a metric value (similar to the procedures in Sections 4.1.3 and 4.1.4 of [RFC9843]) or the link **MAY** be treated as if the metric value is unavailable (e.g., by using a default value). If the Bandwidth metric value is advertised for a link, the PCE **MUST** use the advertised value to compute the path metric in accordance with Sections 4.5.2.1 and 4.5.2.2.

The Path Bandwidth metric value is encoded in IEEE floating point format (see [IEEE.754.2019]).

For use in the PCEP METRIC object, the 24-bit unsigned integer delay value is converted to a 32-bit IEEE floating point value. This conversion follows the procedure specified in [IEEE.754.2019].

4.5.2.1. P2P Path Bandwidth Metric

The Path Bandwidth metric type of the METRIC object in PCEP represents the sum of the Bandwidth metric of all links along a P2P path.

Note: The link Bandwidth metric utilized in the formula may be the original metric advertised on the link, which may have a value inversely proportional to the link capacity.

- A Bandwidth metric of link L is denoted by $B(L)$.
- A Path Bandwidth metric for the P2P path $P = \text{Sum} \{B(L_{pi}), (i=1\dots K)\}$.

4.5.2.2. P2MP Path Bandwidth Metric

The Bandwidth metric type of the METRIC object in PCEP encodes the Path Bandwidth metric for the destination that observes the worst Bandwidth metric among all destinations of the P2MP tree.

- The P2P Bandwidth metric of the path to destination $Dest_j$ is denoted by $BM(Dest_j)$.
- The P2MP Path Bandwidth metric for the P2MP tree $T = \text{Maximum}\{BM(Dest_j), (j=1\dots M)\}$.

4.5.3. User-Defined Metric

Section 2 of [RFC9843] defined a new metric type range for "User-defined metric", which may be advertised in their link metric advertisements. These are user defined and can be assigned by an operator for local use.

User-defined metric values are encoded using the IEEE floating point format (see [IEEE.754.2019]).

For use in the PCEP METRIC object, the 24-bit unsigned integer delay value is converted to a 32-bit IEEE floating point value. This conversion follows the procedure specified in [IEEE.754.2019].

The metric type range was chosen to allow mapping with values assigned in the "IGP Metric-Type" registry. For example, the User-defined metric type 130 of the METRIC object in PCEP can represent the sum of the User-defined metric 130 of all links along a P2P path.

User-defined metrics are equally applicable to P2P and P2MP paths.

5. Operation

The PCEP extensions defined in Sections 5.1 and 5.2 of this document **MUST NOT** be used unless both PCEP speakers have indicated support by setting the S flag in the Path Setup Type sub-TLV corresponding to the PST of the LSP. If this condition is not met, the receiving PCEP speaker **MUST** respond with a PCErr message with Error-Type 19 (Invalid Operation) and Error-value 33 (Attempted use of SR-Algorithm without advertised capability).

The SR-Algorithm used in this document refers to a complete range of SR-Algorithm values (0-255) if a specific section does not specify otherwise. Valid SR-Algorithm values are defined in the "IGP Algorithm Types" registry of the "Interior Gateway Protocol (IGP) Parameters" registry group. Refer to Section 3.1.1 of [RFC8402] and [RFC9256] for the definition of SR-Algorithm in SR. [RFC8665] and [RFC8667] describe the use of the SR-Algorithm in IGP. Note that some RFCs refer to SR-Algorithm with different names, for example, "Prefix-SID Algorithm" and "SR Algorithm".

5.1. ERO and RRO Subobjects

If a PCC receives the Algorithm field in the ERO subobject within PCInitiate, PCUpd, or PCRep messages and the path received from those messages is being included in the ERO of PCRpt message, then the PCC **MUST** include the Algorithm field in the encoded subobjects with the received SR-Algorithm value.

As per [RFC8664], the format of the SR-RRO subobject is the same as that of the SR-ERO subobject but without the L flag; therefore, the SR-RRO subobject may also carry the A flag and Algorithm field in the Subobject Extension Block. Similarly, as per [RFC9603], the format of the SRv6-RRO subobject is the same as that of the SRv6-ERO subobject but without the L flag; therefore, the SRv6-RRO subobject may also carry the A flag and Algorithm field.

5.1.1. SR-ERO

A PCEP speaker **MAY** set the A flag and include the Algorithm field as part of the Subobject Extension Block in an SR-ERO subobject if the S flag has been advertised in the SR-PCE-CAPABILITY sub-TLV by both PCEP speakers.

If the PCEP peer receives an SR-ERO subobject with the A flag set but the S flag was not advertised in SR-PCE-CAPABILITY sub-TLV, then it **MUST** consider the entire ERO as invalid, as described in Section 5.2.1 of [RFC8664].

The Subobject Extension Block field in the SR-ERO subobject **MUST** be included after the optional SID, NAI, or SID structure, and the length of the SR-ERO subobject **MUST** be increased by the size of the Subobject Extension Block, as determined by the set of SEBFs.

If the length and the A flag are not consistent, as specified in Section 4.2, the PCEP peer **MUST** consider the entire ERO invalid and **MUST** send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-value = 11 ("Malformed object").

If the SID value is absent (S flag is set to 1), the NAI value is present (F flag is set to 0), and the Algorithm field is set (the A flag is set to 1), the PCC is responsible for choosing the SRv6-SID value based on values specified in the NAI and Algorithm fields. If the PCC cannot find a SID index in the SR-DB, it **MUST** send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-value = 14 ("Unknown SID").

5.1.2. SRv6-ERO

A PCEP speaker **MAY** set the A flag and include the Algorithm field in an SRv6-ERO subobject if the S flag has been advertised in SRv6-PCE-CAPABILITY sub-TLV by both PCEP speakers.

If the PCEP peer receives an SRv6-ERO subobject with the A flag set or with the SR-Algorithm included, but the S flag was not advertised in SRv6-PCE-CAPABILITY sub-TLV, then it **MUST** consider the entire ERO as invalid, as described in Section 5.2.1 of [RFC8664].

The Algorithm field in the SRv6-ERO subobject **MUST** be included in the position specified in Section 4.3; the length of the SRv6-ERO subobject is not impacted by the inclusion of the Algorithm field.

If the SRv6-SID value is absent (S flag is set to 1), the NAI value is present (F flag is n), and the Algorithm field is set (the A flag is set to 1), the PCC is responsible for choosing the SRv6-SID value based on values specified in the NAI and Algorithm fields. If the PCC cannot find a SID index in the SR-DB, it **MUST** send a PCErr message with Error-Type = 10 ("Reception of an invalid object") and Error-value = 14 ("Unknown SID").

5.2. SR-Algorithm Constraint

To signal a specific SR-Algorithm constraint to the PCE, the headend **MUST** encode the SR-Algorithm TLV inside the LSPA object.

If a PCC receives an LSPA object with the SR-Algorithm TLV as part of PCInitiate or PCUpd messages, then it **MUST** include an LSPA object with the SR-Algorithm TLV in a PCRpt message as part of intended-attribute-list.

If a PCE receives an LSPA object with the SR-Algorithm TLV in PCRpt or PCReq, then it **MUST** include the LSPA object with the SR-Algorithm TLV in a PCUpd message, or a PCRep message in case of an unsuccessful path computation based on rules described in [Section 7.11](#) of [\[RFC5440\]](#).

A PCEP peer that did not advertise the S flag in the Path Setup Type sub-TLV corresponding to the LSP's PST **MUST** ignore the SR-Algorithm TLV on receipt.

The PCE **MUST NOT** use Prefix SIDs associated with an SR-Algorithm other than the one specified in the SR-Algorithm constraint. If a protected Adjacency SID is used without an associated SR-Algorithm, there is a risk that the backup path may fail to forward traffic over parts of the topology that are not included in the specified SR-Algorithm. Consequently, it is **NOT RECOMMENDED** to use protected Adjacency SIDs without an explicitly specified SR-Algorithm. If an Adjacency SID has an associated SR-Algorithm, the PCE **MUST** ensure that the SR-Algorithm matches the one specified in the SR-Algorithm constraint.

Other SID types, such as BSIDs, are allowed. Furthermore, the inclusion of a path BSID from another policy is permitted only if the path associated with that policy fully satisfies all the constraints of the current path computation.

The specified SR-Algorithm constraint is applied to the end-to-end SR Policy path. Using different SR-Algorithm constraints or using winning FAD with different optimization metrics or constraints for the same SR-Algorithm in each domain or part of the topology in single path computation is out of the scope of this document.

If the PCE is unable to find a path with the given SR-Algorithm constraint, it does not support a combination of specified constraints, or if the FAD contains constraints, optimization metrics, or other attributes, which the PCE does not support or recognize, it **MUST** use an empty ERO in PCInitiate for LSP instantiation, a PCUpd message if an update is required, or a NO-PATH object in PCRep to indicate that it was not able to find the valid path.

If the Algorithm field value is in the range 128-255, the PCE **MUST** perform path computation according to the Flexible Algorithm procedures outlined in [Section 5.2.2](#). Otherwise, the PCE **MUST** adhere to the path computation procedures with SID filtering as defined in [Section 5.2.1](#).

If the NO-PATH object is included in PCRep, then the PCE **MAY** include the SR-Algorithm TLV to indicate constraint, which cannot be satisfied as described in [Section 7.5](#) of [\[RFC5440\]](#).

SR-Algorithm does not replace the objective function defined in [\[RFC5541\]](#).

5.2.1. Path Computation for SR-Algorithms 0-127

The SR-Algorithm constraint acts as a filter, restricting which SIDs may be used as a result of the path computation function. Path computation is done based on optimization metric type and constraints specified in the PCEP message received from the PCC.

The mechanism described in this section is applicable only to SR-Algorithm values in the range 0-127. It is not applicable to Flexible Algorithms (range 128-255), which are handled as described in [Section 5.2.2](#). Within the 0-127 range, currently defined algorithms are 0 (Shortest Path First (SPF)) and 1 (Strict-SPF), as introduced in [Section 3.1.1](#) of [\[RFC8402\]](#). Future algorithms defined within this range that do not require explicit PCEP extensions beyond the SR-Algorithm TLV may also utilize this SID filtering approach. If a PCE implementation receives a request with an SR-Algorithm value in the 0-127 range that it does not support for path computation, it **MUST** reject the PCEP message and send a PCerr message with Error-Type 19 (Invalid Operation) and Error-value 34 (Unsupported combination of constraints).

5.2.2. Path Computation for Flexible Algorithms

This section is applicable only to the Flexible Algorithms range of SR-Algorithm values. The PCE performs Flexible Algorithm path computation based on topology information stored in its TED [\[RFC5440\]](#). The TED is expected to be populated with necessary information, including Flexible Algorithm Definitions (FADs), node participation, and ASLA-specific link attributes, through standard mechanisms, such as IGP with Traffic Engineering extensions or BGP - Link State (BGP-LS) [\[RFC9552\]](#).

The PCE must follow the IGP Flexible Algorithm path computation logic as described in [\[RFC9350\]](#). This includes performing the FAD selection as described in [Section 5.3](#) of [\[RFC9350\]](#) and other sections, determining the topology associated with specific a Flexible Algorithm based on the FAD, the node participation ([Section 11](#) of [\[RFC9350\]](#)), using ASLA-specific link attributes ([Section 12](#) of [\[RFC9350\]](#)), and applying other rules for Flexible Algorithm path calculation ([Section 13](#) of [\[RFC9350\]](#)). While [\[RFC9350\]](#) defines the base procedures for IGP Flexible Algorithms, these procedures are further extended by other documents, such as [\[RFC9843\]](#); a PCE implementation may need to support these IGP extensions to allow use of specific constraints in FAD. [\[RFC9917\]](#) created an IANA registry called "IGP Flex-Algorithm Path Computation Rules" <<https://www.iana.org/assignments/igp-parameters>> within the "Interior Gateway Protocol (IGP) Parameters" registry group with the ordered set of rules that **MUST** be used to prune links from the topology during the Flexible Algorithm path computation.

The PCE **MUST** optimize the computed path based on the metric type specified in the FAD. The optimization metric type included in PCEP messages from the PCC **MUST** be ignored. The PCE **MUST** use the metric type from the FAD in messages sent to the PCC unless that metric type is not defined in PCEP or not supported by the PCEP peer. It is allowed to use SID types other than Prefix SID (e.g., Adjacency or BSID) but only from nodes participating in the specified SR-Algorithm.

There are corresponding metric types in PCEP for IGP and TE metrics from FAD introduced in [RFC9350], but there were no corresponding metric types defined for "Min Unidirectional Link Delay" from [RFC9350] and "Bandwidth metric" and "User-defined metric" from [RFC9843]. Section 4.5 of this document introduces them. Note that the defined "Path Bandwidth metric" is accumulative and is different from the BANDWIDTH object defined in [RFC5440].

The PCE **MUST** use the constraints specified in the FAD and also constraints (except optimization metric type) directly included in PCEP messages from the PCC. The PCE implementation **MAY** decide to ignore specific constraints received from the PCC based on existing processing rules for PCEP objects and TLVs, e.g., the P flag described in Section 7.2 of [RFC5440] and processing rules described in [RFC9753]. If the PCE does not support a specified combination of constraints, it **MUST** fail path computation and respond with a PCEP message with a PCInitiate or PCUpd message with an empty ERO or PCRep with NO-PATH object. The PCC **MUST NOT** include constraints from the FAD in the PCEP message sent to the PCE, as it can result in undesired behavior in various cases. The PCE **SHOULD NOT** include constraints from the FAD in PCEP messages sent to the PCC.

The combinations of the constraints specified in the FAD and constraints directly included in PCEP messages from the PCC may decrease the chance that Flexible-Algorithm-specific Prefix SIDs represent an optimal path while satisfying all specified constraints; as a result, a longer SID list may be required for the computed path. Adding more constraints on top of the FAD requires complex path computation and may reduce the benefit of this scheme.

5.3. Metric Types

All the rules of processing the METRIC object as explained in [RFC5440] and [RFC8233] are applicable to the metric types defined in this document.

6. Manageability Considerations

All manageability requirements and considerations listed in [RFC5440], [RFC8231], [RFC8664], and [RFC9603] apply to the PCEP extensions defined in this document. In addition, the requirements and considerations listed in this section apply.

6.1. Control of Function and Policy

A PCE or PCC implementation **MAY** allow the capability of supporting the PCEP extensions introduced in this document to be enabled or disabled as part of the global configuration. By default, this capability **SHOULD** be enabled.

6.2. Information and Data Models

An implementation **SHOULD** allow the operator to view the capability defined in this document. Sections 4.1 and 4.1.1 of [RFC9826] should be extended to include the capabilities introduced in Sections 4.1.1 and 4.1.2 for the PCEP peer.

6.3. Liveness Detection and Monitoring

This document does not define any new mechanism that impacts the liveness detection and monitoring of PCEP.

6.4. Verify Correct Operations

An implementation **SHOULD** also allow the operator to view FADs, which may be used in Flexible Algorithm path computation as defined in [Section 5.2.2](#).

An implementation **SHOULD** allow the operator to view nodes participating in the specified SR-Algorithm.

6.5. Requirements on Other Protocols and Functional Components

This document does not put new requirements but relies on the necessary IGP extensions.

6.6. Impact on Network Operations

This document inherits considerations from documents describing IGP Flexible Algorithm -- for example, [\[RFC9350\]](#) and [\[RFC9843\]](#).

7. Operational Considerations

This document inherits operational considerations from documents describing IGP Flexible Algorithm -- for example, [\[RFC9350\]](#) and [\[RFC9843\]](#).

8. Security Considerations

The security considerations described in [\[RFC5440\]](#), [\[RFC8231\]](#), [\[RFC8253\]](#), [\[RFC8281\]](#), [\[RFC8664\]](#), [\[RFC9603\]](#), and [\[RFC9350\]](#) apply to the extensions described in this document as well.

Note that this specification introduces the possibility of computing paths by the PCE based on Flexible-Algorithm-related topology attributes and based on the metric type and constraints from the FAD. This creates additional vulnerabilities, which are already described for the path computation done by IGP, like those described in the Security Considerations section of [\[RFC9350\]](#) but which are also applicable to path computation done by the PCE. Hence, securing the PCEP session using Transport Layer Security (TLS) [\[RFC8253\]](#) [\[RFC9916\]](#) is **RECOMMENDED** as per the recommendations and best current practices described in [\[RFC9325\]](#).

9. IANA Considerations

9.1. SR Capability Flag

IANA maintains a registry named "SR Capability Flag Field" within the "Path Computation Element Protocol (PCEP) Numbers" registry group to manage the Flags field of the SR-PCE-CAPABILITY sub-TLV. IANA has registered the following:

Bit	Description	Reference
5	SR-Algorithm Capability	RFC 9933

Table 1

9.2. SRv6 PCE Capability Flag

IANA maintains a registry named "SRv6 Capability Flag Field" within the "Path Computation Element Protocol (PCEP) Numbers" registry group to manage the Flags field of SRv6-PCE-CAPABILITY sub-TLV. IANA has registered the following:

Bit	Description	Reference
13	SR-Algorithm Capability	RFC 9933

Table 2

9.3. SR-ERO Flag

IANA maintains a registry named "SR-ERO Flag Field" within the "Path Computation Element Protocol (PCEP) Numbers" registry group to manage the Flags field of the SR-ERO Subobject. IANA has registered the following:

Bit	Description	Reference
7	SR-Algorithm Flag (A)	RFC 9933

Table 3

9.4. SRv6-ERO Flag

IANA maintains a registry named "SRv6-ERO Flag Field" within the "Path Computation Element Protocol (PCEP) Numbers" registry group to manage the Flags field of the SRv6-ERO subobject. IANA has registered the following:

Bit	Description	Reference
7	SR-Algorithm Flag (A)	RFC 9933

Table 4

9.5. PCEP TLV Types

IANA maintains a registry named "PCEP TLV Type Indicators" within the "Path Computation Element Protocol (PCEP) Numbers" registry group. IANA has registered the following TLV type for the new LSPA TLV specified in this document.

Value	Description	Reference
66	SR-Algorithm	RFC 9933

Table 5

9.6. Metric Types

IANA maintains a registry named "METRIC Object T Field" within the "Path Computation Element Protocol (PCEP) Numbers" registry group. IANA has registered these codepoints as follows:

Value	Description	Reference
22	Path Min Delay metric	RFC 9933
23	P2MP Path Min Delay metric	RFC 9933
24	Path Bandwidth metric	RFC 9933
25	P2MP Path Bandwidth metric	RFC 9933
128-255	User-defined metric	RFC 9933

Table 6

9.7. PCEP-Error Object

IANA has registered the following Error-Types and Error-values within the "PCEP-ERROR Object Error Types and Values" registry of the "Path Computation Element Protocol (PCEP) Numbers" registry group.

Error-Type	Meaning	Error-value	Reference
19	Invalid Operation	33: Attempted use of SR-Algorithm without advertised capability	RFC 9933

Error-Type	Meaning	Error-value	Reference
		34: Unsupported combination of constraints	RFC 9933

Table 7

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